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SUMMER – 19 EXAMINATION

Subject Name: Data Communication Network Model Answer

Subject Code: 22414

Important Instructions to examiners:

- 1) The answers should be examined by key words and not as word-to-word as given in the model answer scheme.
- 2) The model answer and the answer written by candidate may vary but the examiner may try to assess the understanding level of the candidate.
- 3) The language errors such as grammatical, spelling errors should not be given more Importance (Not applicable for subject English and Communication Skills.
- 4) While assessing figures, examiner may give credit for principal components indicated in the figure. The figures drawn by candidate and model answer may vary. The examiner may give credit for any equivalent figure drawn.
- 5) Credits may be given step wise for numerical problems. In some cases, the assumed constant values may vary and there may be some difference in the candidate's answers and model answer.
- 6) In case of some questions credit may be given by judgement on part of examiner of relevant answer based on candidate's understanding.
- 7) For programming language papers, credit may be given to any other program based on equivalent concept.

| Q. | Sub | Answer | Marking |
|-----|------------|--|------------|
| No. | Q . | | Scheme |
| | N. | | |
| Q.1 | | Attempt any five of the following: | 10 M |
| | | | |
| | a | Define Computer Network and state its types. | 2 M |
| | Ans | Definition: | 1 M |
| | | A computer network is a group of computer systems and other computing hardware | definition |
| | | devices that are linked together through communication channels to facilitate | , 1M for |
| | | communication and resource sharing among a wide range of users. | types |
| | | | |
| | | Types of Computer Networks: | |
| | | | |
| | | • Local Area Networks (LAN) | |
| | | • Personal Area Networks (PAN) | |
| | | Home Area Networks (HAN) | |
| | | • Wide Area Networks (WAN) | |
| | | Metropolitan Area Networks (MAN) | |
| | | • The Internet | |
| | | | |
| | b | State various Computer Network applications | 2 M |
| | Ans | Computer Network Applications: | Any |
| | | 1. File Sharing | Four- 1/2 |



| | 2. Printe | r Sharing | | Γ | M each |
|-----|--|---------------------------------|-------------------------------|-----|----------|
| | 3. Application Services | | | | |
| | 4. E-mail Services | | | | |
| | 5. Remo | te access | | | |
| | 6. Interne | et & Intranet | | | |
| | | | | | |
| С | List any | four Unguided Transmission Me | edia. | | 2M |
| Ans | Unguided | d Media or Wireless media: | | 1/2 | M each |
| | (a) Radio | wave | | | |
| | (b) Micro | owave | | | |
| | (c) infrar | ed | | | |
| | (d) Satell | ite | | | |
| d | State typ | oes of Errors | | | |
| Ans | Content l | Error | | 1 | M each |
| | Flow Integrity error | | | | |
| e | List IEEE 802 X standards for networks | | | | 2M |
| Ans | 1. 802.3: | Ethernet | | | 1/2 M |
| | 2 002 4 | | | | each |
| | 2. 802.4 | Token Bus | | | |
| | 3. 802.5:Token Ring | | | | |
| | 4. 802.1 | 1:Wi Fi(Wireless Fidelity) | | | |
| f | Compar | e Router and Repeater. | | | 2M |
| Ans | | Router | Repeater | | any 2 |
| | | | | p | points 1 |
| | | A router is a device like a | Repeater regenerates the | I | M each |
| | | switch that routes data packets | network before the signal | | |
| | | based on their IP addresses. | becomes too weak or | | |
| | | | corrupted so as to extend the | | |
| | | | length to which the signal | | |
| | | | can be transmitted over the | | |
| | | | same network. | | |
| | | Router is mainly a Network | A repeater operates at the | | |
| | | Layer device. | physical layer. | | |
| | | | 1 | | |



| | g Ans | State functions of Network layer Functions of network layer: 1. logical addressing 2. Routing. 3. Congestion control 4. Accounting and billing | 2M 1/2M each | | |
|----|----------|---|---------------------------|--|--|
| | | 4. Accounting and billing | | | |
| | | 6. Source host to destination host error free delivery of packet | | | |
| | | | | | |
| | | | | | |
| Q2 | | Attempt any THREE of the following : | 12 M | | |
| | a | Classify the network based on geographical area and transmission technology | 4 M | | |
| | Ans | Classification of networks based on geography: | 2 M for | | |
| | | | geographi | | |
| | | LAN - Local Area Network | cal area | | |
| | | MAN - Metropolitan Area Network | and 2 M | | |
| | | What we | tor | | |
| | | WAN - Wide Area Network | ion | | |
| | | | technolog | | |
| | | CAN - Campus Area Network | v. | | |
| | | PAN - Personal Area Network | Explanati | | |
| | | LAN: LAN is local area network. LAN is privately-owned networks covering a small geographic area(less than 1 km), like a home, office, building or group of buildings. LAN transmits data with a speed of several megabits per second. | | | |
| | | MAN: A Metropolitan Area Network (MAN) is a large computer network that spans a metropolitan area or campus. 2. A MAN typically covers an area up to 10 kms (city). The best example of MAN is the cable Television network, available in many cities. | | | |



| | WAN: WAN is wide area network. WAN is a long-distance communication network that covers a wide geographic area, such as state or country. The most common example is internet. | |
|-----|--|--|
| | The transmission technology can be categorized broadly into two types: | |
| | 1. Broadcast networks | |
| | Broadcast networks have a single communication channel that is shared or used by all the machines on the network. Short messages called packets sent by any machine are received by all the others. Broadcast systems generally use a special code in the address field for addressing a packet to all the concerned computers. This mode of operation is called broadcasting. | |
| | 2. Point-to-point networks | |
| | Point to point networks consists of many connections between individual pairs of | |
| | machines. To go from the source to the destination a packet on these types of | |
| | network may have to go through intermediate computers before they reach the | |
| | desired computer. | |
| h | Draw structural diagram of fiber ontic cable and write its functions | 4 M |
| Ans | Druw structural diagram of more optic cubic and write its functions | 2 M for |
| | Black polyurethane outer jacket | diagram and 2 M for functions |
| | Fig. Structural diagram for Fibre Optic Cable | |
| | Fig. Structural diagram for Fibre Optic Cable Functions of Optical Cable: | |
| | Fig. Structural diagram for Fibre Optic Cable Functions of Optical Cable: 1. Single-mode fibers - Used to transmit one signal per fiber (used in telephones and cable TV) | |
| | Fig. Structural diagram for Fibre Optic Cable Functions of Optical Cable: 1. Single-mode fibers - Used to transmit one signal per fiber (used in telephones and cable TV) 2. Multi-mode fibers - Used to transmit many signals per fiber (used in computer | |



| | networks, local area networks) | |
|-----|---|--|
| с | Describe various IEEE standards for network topologies. | 4 M |
| Ans | A set of network standards developed by the IEEE. They include: | 1 Mark for 1 |
| | IEEE 802.1: Standards related to network management. IEEE 802.2: General standard for the data link layer in the OSI Reference Model. The IEEE divides this layer into two sublayers the logical link control (LLC) layer and the media access control (MAC) layer. The MAC layer varies for different network types and is defined by standards IEEE 802.3 through IEEE 802.5. IEEE 802.3: Defines the MAC layer for bus networks that use CSMA/CD. This is the basis of the Ethernet standard. EEE 802.4: Defines the MAC layer for bus networks that use a token-passing mechanism (token bus networks). IEEE 802.5: Defines the MAC layer for token-ring networks. IEEE 802.6: Standard for Metropolitan Area Networks (MANs). IEEE 802.11 Wireless Network Standards: 802.11 is the collection of standards setup for wireless networking. | each |
| d | Draw and explain layered architecture of OSI model. | 4 M |
| Ans | OSI model (Open System Interconnection) model was developed by ISO (international standard organization) which provides way to understand how internetwork operates. It gives guidelines for creating network standard. OSI model has 7 layers as shown in the figure. Application Layer, Presentation Layer ,Session Layer ,Transport Layer ,Network Layer ,Data link Layer and Physical Layer | 1 M diagram and 3 M explanati on |
| | Physical (Layer 1) OSI Model, Layer 1 conveys the bit stream - electrical impulse, light or radio signal — through the network at the electrical and mechanical level. It provides the hardware means of sending and receiving data on a carrier, including defining cables, cards and physical aspects. | |
| | Data Link (Layer 2) At OSI Model, Layer 2, data packets are encoded and decoded into bits. It furnishes transmission protocol knowledge and management and handles errors in the physical layer, flow control and frame synchronization. The data link layer is divided into two sub layers: The Media Access Control (MAC) layer and the Logical Link Control (LLC) layer. The MAC sub layer controls how a computer on the network gains access to the data and permission to transmit it. The LLC layer controls frame synchronization, flow control and error | |







| | | application services for file. | | |
|----|-----|---|--|------------------------|
| | | | | |
| 03 | | Attempt any THREE of the following | o • | 12 M |
| | a | What advantages does TDM have ov | er FDM in a circuit switched network? | 4 M |
| | Ans | In TDM, each signal uses all of the ba | andwidth some of the time, while for FDM, | consider |
| | | each signal uses a small portion of the | bandwidth all of the time. | 4 points for 4 M |
| | | TDM uses the entire frequency range | but dynamically allocates time, certain jobs | |
| | | might require less or more time, which | n TDM can offer but FDM is unable to as it | |
| | | cannot change the width of the allocate | d frequency. | |
| | | TDM provides much better flexibility c | compared to FDM. | |
| | | TDM offers efficient utilization of band | dwidth | |
| | | Low interference of signal and minimiz | zes cross talk | |
| | h | Compare Analog and Digital signal | | 4 M |
| | Ans | Analog signal | Digital signal | 1 M for |
| | | An analog signal is a continuous wave that changes over a time period. | A digital signal is a discrete wave that carries information in binary form. | each differenc e |
| | | An analog signal is represented by a sine wave. | A digital signal is represented by square waves. | any 4 valid |
| | | Analog signal has no fixed range. | Digital signal has a finite numbers i.e. 0 and 1. | points |
| | | An analog signal is described by the amplitude, period or frequency, and phase. | A digital signal is described by bit rate and bit intervals. | |
| | | An analog signal is more prone to distortion. | A digital signal is less prone to distortion. | |
| | | An analog signal transmits data in the form of a wave. | A digital signal carries data in the binary form i.e. 0 and 1. | |
| | | | | |
| | C | With suitable diagram describe | | 2M - 4 |
| | Ans | Star topology is a network topology v | where each individual piece of a network is | topology- 1M for |



attached to a central node (often called a hub or switch). The attachment of these network pieces to the central component is visually represented in a form similar to a star. diagram

The hub and hosts, and the transmission lines between them, form a graph with the topology of a star. Data on a star network passes through the hub before continuing to its destination. The hub manages and controls all functions of the network. It also acts as a repeater for the data flow.



Fig a: Star topology

The star network is one of the most common computer network topologies.

(ii)**RING Topology**

A ring network is a network topology in which each node connects to exactly two other nodes, forming a single continuous pathway for signals through each node - a ring.

Data travels from node to node, with each node along the way handling every packet.



and 1 mark for descriptio n ,2M ring topology-1 M for diagram and 1 Mark for descriptio

n



| | | Fig b: Ring Topology | |
|----|-----|--|---------------------------------|
| | d | Ring topology refers to a specific kind of network setup in which devices are connected in a ring and pass information to or from each other according to their adjacent proximity in the ring structure. This type of topology is highly efficient and handles heavier loads better than bus topology. Describe the major functions of network layer in TCP/IP protocol suite | |
| | Ans | Internetworking: This is the main duty of network layer. It provides the logical connection between different types of networks. Addressing: Addressing is necessary to identify each device on the internet uniquely. This is similar to telephone system. The address used in the network layer should uniquely and universally define the connection of a computer. Routing: In a network, there are multiple roots available from a source to a destination and one of them is to be chosen. The network layer decides the root to be taken. This is called as routing. Packetizing: The network layer encapsulates the packets received from upper layer protocol and makes new packets. This is called as packetizing. It is done by a network layer protocol called IP (Internetworking Protocol). | 1 M for each function |
| 04 | | Attempt any Five of the following: | 12 M |
| 4 | а | Draw and describe architecture for network using tree topology for an office | 4 M |
| | | in 3-storeys building. | |
| | Ans | A tree topology is a special type of structure in which many connected elements are arranged like the branches of a tree Here in the diagram the main switch is connected with three separate switches. For each floor separate switch is connected with multiple terminals. | Explain 1M ,Diagram 3M |







| | Class A range for first byte is 0-127. | Class A type of IP addresses have First | byte M |
|-------|---|--|-------------------------------|
| ans | Class A: | | Explain 4 |
| d | Describe types of IP address classes. | | 4 M |
| | | | |
| | Interference is high | Interference is Low or negligible | |
| | | analog signals | |
| | Used with Analog signals | Used with both Digital signals and | |
| | Frequency is shared | Times scale is shared | |
| | FDM divides the channel into two or more frequency ranges that do not overlap | TDM divides and allocates certain time periods to each channel in an alternating manner | differenc e |
| ans | Frequency Division Multiplexing | Time division Multiplexing | each |
| c | Differentiate between FDM and TDM | 1 Time disisten Markinlasia a | 4 M |
| | • Another function of data link 1 devices are attached to the sam which device has control over the | ayer is access control. When two or more link, data link layer protocols determe link at any given time. | more mine |
| | • It also provide error control moduplicate, or lost frame, thus add | echanism to detect & retransmit dama ling reliability to physical layer. | nged, |
| | • It provides flow control mechan data at the speed that the receive | ism to ensure that sender is not sending or cannot process. | g the |
| | • Physical addresses of source & frame. | & destination machines are added to | each |
| | • It then provides the addressin frame. | ng information by adding header to | each |
| | • Data link layer receives the dat manageable units called frames. | ta from the network layer & divide it | into |
| | Functions of Data Link Layer | | |
| | • It also includes the mechanica different network interfaces | al, electrical and timing specifications | s for |
| | • Physical layer therefore encomp in different types of Wired/Wire links connecting them | basses the set of all protocols/standards eless interfaces and the telecommunication | used functions ation each) |



| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0 Class C: Class C range for first byte is 192-223. This class has first three bytes referring to network with starting bits as 110 and last byte signifies Host ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.255.0 The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 240-255 .Class E is reserved for multicasting and its starting bits are 1110 Class E: Class E range for first byte is 240-255 .Class E is reserved for future use and its starting bits are 1111 Elass D NET ID HOST ID HOST ID Class B NET ID Host ID HOST ID Fig : IP address classes | e | Design suitable netw | work layout for an organization with five department | 4 M |
|---|---|--|---|-----|
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. Class C: Class C range for first byte is 192-223. This class has first three bytes referring to network with starting bits as 110 and last byte signifies Host ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.255.0. The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 224-239 Class D is used for multicasting and its starting bits are 1110 Class E: Class E range for first byte is 240-255 .Class E is reserved for future use and its starting bits are 1111 Class D NET ID HOST ID HOST ID Class E NET ID HOST | | | Fig : IP address classes | |
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. Class C: Class C range for first byte is 192-223. This class has first three bytes referring to network with starting bits as 110 and last byte signifies Host ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.255.0 The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 224-239 Class D is used for multicasting and its starting bits are 1110 Class E: Class E range for first byte is 240-255. Class E is reserved for future use and its starting bits are 1111 Class A NET ID Host ID HOST ID Class C NET ID Host ID HOST ID Class C NET ID Host ID HOST ID Class C NET ID Host ID HOST ID Class D MULTICAST ADDRESS <th></th> <th>Class E</th> <th>RESERVED</th> <th></th> | | Class E | RESERVED | |
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. Class C: Class C range for first byte is 192-223. This class has first three bytes referring to network with starting bits as 110 and last byte signifies Host ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.255.0 The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 224-239 Class D is used for multicasting and its starting bits are 1110 Class E: Class E range for first byte is 240-255 .Class E is reserved for future use and its starting bits are 1111 | | Class D | MOLTICAST ADDRESS | |
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. Class C: Class C range for first byte is 192-223. This class has first three bytes referring to network with starting bits as 110 and last byte signifies Host ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.050. The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 240-255. Class E is reserved for future use and its starting bits are 1111 | | Class C | | |
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. Class C: Class C range for first byte is 192-223. This class has first three bytes referring to networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.255.0. The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 224-239 Class D is used for multicasting and its starting bits are 1110 Class A: Payle 1 Byte 2 Byte 3 Byte 4 HOST ID | | Class B | NET ID HOST ID | |
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. Class C: Class C range for first byte is 192-223. This class has first three bytes referring to network with starting bits as 110 and last byte signifies Host ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.05.0 The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 240-255 .Class D is used for multicasting and its starting bits are 1111 | | Class A | NET ID HOST ID | |
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. Class C: Class C range for first byte is 192-223. This class has first three bytes referring to network with starting bits as 110 and last byte signifies Host ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.0.0. The network ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.0.0. The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 224-239 Class D is used for multicasting and its starting bits are 1110 Class E: Class E range for first byte is 240-255 .Class E is reserved for future use and its starting bits are 1111 | | | Byte 1 × Byte 2 × Byte 3 × Byte 4 | |
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. Class C: Class C range for first byte is 192-223. This class has first three bytes referring to network with starting bits as 110 and last byte signifies Host ID. Here, number of networks is more when compared to number of hosts in each network. The default subnet masks for class C is 255.255.255.0 The network IP addresses for these range from 192.0.0.0 to 223.0.0.0. Class D: Class D range for first byte is 224-239 Class D is used for multicasting and its starting bits are 1110 | | Class E : Class E ra and its starting bits | nge for first byte is 240-255 .Class E is reserved for future use are 1111 | |
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| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. Class B: Class B range for first byte is 128-191. This type has first two bytes specifying network ID with starting two bits as 10 and last two bytes referring to host ID. The default subnet masks for class B is 255.255.0.0. Network addresses for these ranges from 128.0.0.0 to 191.0.0.0. | | Class C : Class C r referring to network number of network The default subnet | ange for first byte is 192-223. This class has first three bytes a with starting bits as 110 and last byte signifies Host ID. Here, is is more when compared to number of hosts in each network. masks for class C is 255.255.255.0 The network IP addresses | |
| consisting of Network address with first bit as 0 and the next 3 bytes with host id. Hence, number of hosts are more when compared to number of networks. The default subnet masks for class A networks is 255.0.0.0. Class A networks have their network addresses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by node addresses. | | Class B : Class B a specifying network host ID. The defau for these ranges from | range for first byte is 128-191. This type has first two bytes ID with starting two bits as 10 and last two bytes referring to lt subnet masks for class B is 255.255.0.0. Network addresses m 128.0.0.0 to 191.0.0.0. | |
| | | consisting of Netwo Hence, number of default subnet mast their network addre node addresses. | ork address with first bit as 0 and the next 3 bytes with host id. hosts are more when compared to number of networks. The ks for class A networks is 255.0.0.0. Class A networks have sses from 1.0.0.0 to 126.0.0.0, with the zero's being replaced by | |



| | ans | Internet Fixewalt Dept 2. Dept 3. Fixewalt Dept 4. Dept 5 | ten users each) (Correct dia 4M) Consider any suitable diagram |
|----|-----|--|---|
| 05 | | Attempt any TWO of the following: | 12 M |
| | a | Describe the process of data communication in various modes | 6 M |
| | ans | Transmission mode refers to the mechanism of transferring of data between two | mode |
| | | devices connected over a network. It is also called Communication Mode. These modes direct the direction of flow of information. There are three types of transmission modes. They are: | explanati on 1 M each & diagram 1 M each |
| | | Simplex Mode Half duplex Mode Full duplex Mode a. In Simplex mode, the communication is unidirectional, as on a one-way street. Only one of the two devices on a link can transmit; the other can only receive. The simplex mode can use the entire capacity of the channel to send data in one direction. Keyboards, traditional monitors and printers are examples of simplex devices. | |



| | Direction of data Monitor Mainframe Simplex Mode a. In half-duplex mode, each station can both transmit and receive, but not at the same time. When one device is sending, the other can only receive, and vice versa. The half-duplex mode is used in cases where there is no need for communication in both directions at the same time. The entire capacity of the channel can be utilized for each direction -for example :Walkie-talkies. Image: Direction of data at time 1 Image: Direction of data at time 2 Workstation Direction of data at time 3 Workstation Direction of data at time 4 Workstation Direction of data at time 5 Workstation Direction of data at time 5 One common example of full-duplex communication is the telephone net | |
|------|---|-------------------------------|
| | | |
| D | communication? | |
| ans | Switching is a mechanism by which data/information sent from source towards destination which are not directly connected. Networks have interconnecting devices, which receives data from directly connected sources, stores data, analyse it and then forwards to the next interconnecting device closest to the destination. Switching can be categorized as: Circuit switching | Any six points 1 M each |
| | Packet switching | |



| Circuit switching is preferred over packet switching in voice communication because: | |
|---|-----------|
| because: | |
| | |
| • In circuit switching a dedicated nath is established between sender and | |
| receiver which is maintained for entire duration of conversation | |
| • It provides continuous and guaranteed delivery of date | |
| • It provides continuous and guaranteed derivery of data. | |
| • During the data transfer phase, no addressing is needed. | |
| • Delays are small. | |
| It uses connection oriented service. | |
| Message received in order to the destination | |
| | |
| c Your company has the network id 165.130.0.0. You are responsible for | 6 M |
| creating subnets on the network, and each subnet must provide at least 1000 | |
| host ids. What subnet mask meets the requirement for the minimum number | |
| of host ids and provides the highest number of subnets? | |
| ans The given network id 165.130.0.0 is class B (Range of class B is 128.0.0.0 to Ez | Explanati |
| 191 255 255 255) with subnet mask of 255 255 252 0 creates 62 subnets with 1022 | on 6 M |
| host each | |
| nost cach. | |
| In binary format subnet mask reads: | |
| | |
| 111111111111111111100.00000000. | |
| | |
| To calculate the number of host ids available for each subnet is based on the | |
| number of digits remaining in the network address. | |
| | |
| The number of possible host ids in each subnet ranges from 00000001 through | |
| 11111110. | |
| | |
| So, in the network 165.130.0.0/22, host addresses can range from 165.130.0.1 | |
| through 165.130.254 | |
| | |
| | |
| Q6 Attempt any TWO of the following: | |
| a A system uses CRC on a block of 8 bytes. How many redundant bits are sent | 6 M |
| per block? What is the ratio of useful bits to total bits? | |
| ans CRC is one of the most common and powerful error detecting code which can be De | Descripti |
| describe as follows. The polynomial code also known as CRC with co-efficient of o | on 6 M |
| 0s and 1s. In this method the sender and receiver must agree upon generator *T | The |
| polynomial g(x) in advance. Both the high and low order bits of the generator stu | tudent |
| (divisor) must be 1. To compute the checksum for some frame (data) with m bits, m | nay |
| the frame must be longer than generator polynomial. The idea is to append as | ssume a |
| checksum to the end of frame in such a way that the polynomial represented by the po | olynomi |
| checksum frame is divisible by $g(x)$. When the receiver gets the checksum frame it a | lora |



| | tries dividing it by $g(x)$. If there is remainder there has been a transmission error and zero remainder means no error in the transmission. r is degree of $g(x)$ polynomial. Step by step procedure: 1. Append a string of r zero bits to the lower order end of data word(m) where r is less than the number of bits pre-decided divisor by 1 bit i.e. if divisor = 5 bits then r = 4 zeros. Now data word contains m+r bits 2. Divide the newly generated data unit in step 1 by the divisor. It is module – 2 division 3. The remainder obtained after division is the r bit CRC. 4. This CRC will replace the r zeros appended to the data unit to get the code word to be transmitted. NOTE: The polynomial code for calculation of redundant bits is not given .hence the data given is insufficient for calculating redundant bits and the ratio of useful bits to total bits. | divisor and do the problem. Full marks has to be given even if they explain the method or do the problem with assumpti ons'. |
|-----|---|---|
| b | Describe the process of DHCP server configuration. | 6 M |
| ans | DHCP (Dynamic Host Configuration Protocol) is a client-server protocol that uses DHCP servers and DHCP clients. A DHCP server is a machine that runs a service that can lease out IP addresses and other TCP/IP information to any client that requests them. The DHCP server typically has a pool of IP addresses that it is allowed to distribute to clients, and these clients lease an IP address from the pool for a specific period of time, usually several days. Once the lease is ready to expire, the client contacts the server to arrange for renewal. DHCP clients are client machines that run special DHCP client software enabling them to communicate with DHCP server. | Diagram 2M, Explanati on 4 M |
| | DHCPOFFER DHCPOFFER DHCPACK DHCPACK DHCP client DHCPACK DHCP server DHCP server DHCP clients obtain a DHCP lease for an IP address, a subnet mask, and various DHCP options from DHCP servers in a four-step process: DHCP DISCOVER: The client broadcasts a request for a DHCP server. DHCPOFFER: DHCP servers on the network offer an address to the client. | |
| | DHCPREQUEST : The client broadcasts a request to lease an address from one of | |



| | the offering DHCP servers. | |
|-----|--|--|
| | DHCPACK : The DHCP server that the client responds to acknowledges the client, assigns it any configured DHCP options, and updates its DHCP database. The client then initializes and binds its TCP/IP protocol stack and can begin network communication. | |
| c | What is the MAC protocol used in TOKEN ring LAN's? What happens if the token is lost? | 6 M |
| ans | Token ring local area network (LAN) network is a communication protocol for local area networks.it uses special three-byte frame called a "token" that travels around a logical ring of workstations or servers. This token passing is a channel access method providing fair access for all stations, and eliminating the collision of contention-based access methods. Introduced by IBM in 1984, it was then standardized with protocol IEEE 802.5 and was fairly successful, particularly in the corporate environments, but gradually eclipsed by the later versions of Ethernet. The IEEE 802.5 Token ring technology provides for data transfer rates of either 4 | Descripti on of MAC protocol 4 M, Explanati on of token lost 2 M |
| | or 16 Mbps. | |
| | It works in the following manner: | |
| | 1. Empty information frames are continuously circulated on the ring. | |
| | 2. When a computer has a message to send, it inserts a token in an empty frame (simply changing a 0 to a 1 in the token bit part of the frame) and a message and a destination identifier in the frame. | |
| | 3. The frame is the examined by each successive workstation. If workstation sees that it is the destination of the message, it copies the message from the frame and changes the token back to 0. | |
| | 4. When the frame gets back to originator, it sees that message has been copied and received. | |
| | The Fibre Distributed Data Interface (FDDI) also uses a Token ring protocol. | |
| | If one device does not receive a token within a specified period, it can issue an alarm. The alarm alerts the network administrator to the problem and its location. Then, network administrator generates a new , free token | |
| | OR | |



| \cdot There are two error conditions that could cause the token ring to break down. | |
|--|--|
| One is the lost token in which case there is no token in the ring.Other is the busy token that circulates endlessly. | |
| To overcome these problems, the IEEE 802 standard specifies that one of the stations must be designated as "active monitor". The monitor detects the lost condition using a timer by time-out mechanism and recovers by using a new free token | |