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WINTER - 19 EXAMINATIONS

Subject Name: Digital Communication Systems Model Answer

Subject Code: 22428

Important Instructions to examiners:

- 1) The answers should be examined by key words and not as word-to-word as given in the model answer scheme.
- 2) The model answer and the answer written by candidate may vary but the examiner may try to assess the understanding level of the candidate.
- 3) The language errors such as grammatical, spelling errors should not be given more Importance (Not applicable for subject English and Communication Skills.
- 4) While assessing figures, examiner may give credit for principal components indicated in the figure. The figures drawn by candidate and model answer may vary. The examiner may give credit for any equivalent figure drawn.
- 5) Credits may be given step wise for numerical problems. In some cases, the assumed constant values may vary and there may be some difference in the candidate's answers and model answer.
- 6) In case of some questions credit may be given by judgement on part of examiner of relevant answer based on candidate's understanding.
- 7) For programming language papers, credit may be given to any other program based on equivalent concept.

Q. No.	Sub Q. N.	Answer	Marking Scheme
Q.1		Attempt any FIVE of the following:	10- Total Marks
	a)	Define (i)Bit rate (ii)Baud rate	2M
	Ans:	(i)Bit rate :- Bit rate is simply the number of bits transmitted during one second and is expressed in bits per second (bps). Mathematically bit rate is given by:- $Rb = 1 / T_b$ where Tb is time interval of one bit	1M each
		(ii)Baud rate: - Baud is the unit of symbol rate. Baud rate is the number of symbols transmitted during one second and is expressed in symbols per second or baud. Mathematically, baud rate is the reciprocal of the time of one output signaling element and a Signaling element (symbol) may represent several information bits. Baud rate is expressed as, $Rs=1/T_s$ Where, baud rate = symbol rate (symbols per second) and T _s = time interval of one symbol.	
	b)	State the Hartley's law with mathematical expression.	2M
	Ans:	Hartley's Law / Nyquist Theorem:- Statement: Hartley's Theorem/Law states that the channel capacity of the transmission channel of bandwidth 'B' which carries a signal having 'M' levels in the total absence of noise is given by: $C = 2 B \log 2 M$	
		 where, C – channel capacity (bits/sec) B – channel bandwidth M – number of coding levels (2 or more) In the absence of noise, Hartley's Law shows that greater the number of levels in the coding system, the greater the information rate that can be sent through the channel. 	

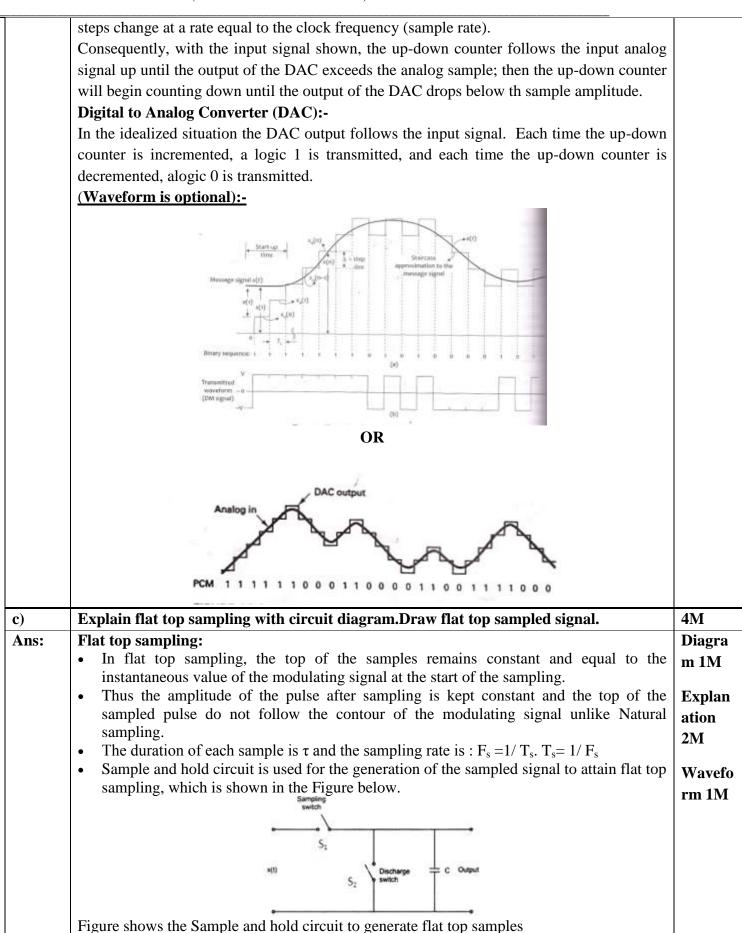


c)	State sampling theorem. Define Nyquist rate.	2M
Ans:	SAMPLING THEOREM:	1M
	Sampling theorem states that a band-limited signal of finite energy having the highest	each
	frequency component f_m Hz can be represented and recovered completely from a set of	
	samples taken at a rate of f_s samples per second provided that $f_s \ge 2f_m$.	
	Here f_s is the sampling frequency. This theorem is also known as the Sampling Theorem for	
	Baseband or Low-pass Signals.	
	Nyquist rate:-	
	Sampling frequency should be equal to or greater than twice the maximum signal	
	frequency $(f_s \ge 2f_m)$	
d)	Classify the modulation techniques.	2M
Ans:	Classification of the modulation techniques:-	2M
	1. Amplitude Shift Keying (ASK)	
	2. Frequency Shift Keying (FSK)	
	3. Phase Shift Keying (PSK)	
e)	State two advantages of WDM technique.	2M
Ans:	ADVANTAGES OF WDM:	Any 2
	1. WDM has enhanced capacity.	1M
	2. WDM can be used for full duplex transmission with a single fiber.	each
	3. It is inherently easier to reconfigure (addition or removal of channels).	
	4. Fiber optic cable networks use optical components which are simpler and more reliable	
	and often less costly than their electronic counterparts	
f)	List the various multiple access techniques.	2M
Ans:	1. Frequency Division Multiple Access (FDMA)	½ M
	2. Time Division Multiple Access (TDMA)	each
	3. Code Division Multiple Access (CDMA)	
	4. Space Division Multiple Access (SDMA)	
g)	Define the concept of spread spectrum.	2M
Ans:	Concept of spread spectrum :-	2M
	Spread-spectrum techniques are methods by which a signal (e.g. an electrical,	
	electromagnetic, or acoustic signal) generated with a particular <u>bandwidth</u> is deliberately	
	spread in the <u>frequency domain</u> , resulting in a signal with a wider <u>bandwidth</u> . OR	
	Spread spectrum systems are intended to provide such secure and reliable communication.	
	In this system the spectrum of the transmitted signals spreaded over a very wide	
	bandwidth. This achieved in these systems by modulating for a second time, an already	
	modulated signal in such a way as to spread the power of the transmitted spread spectrum	
	signal over a very large bandwidth.	
	Attempt any THREE of the following:	12-
	Attempt any THREE of the following:	Total
		Mark
a)	State the advantages and disadvantages of digital communication system.	4 M

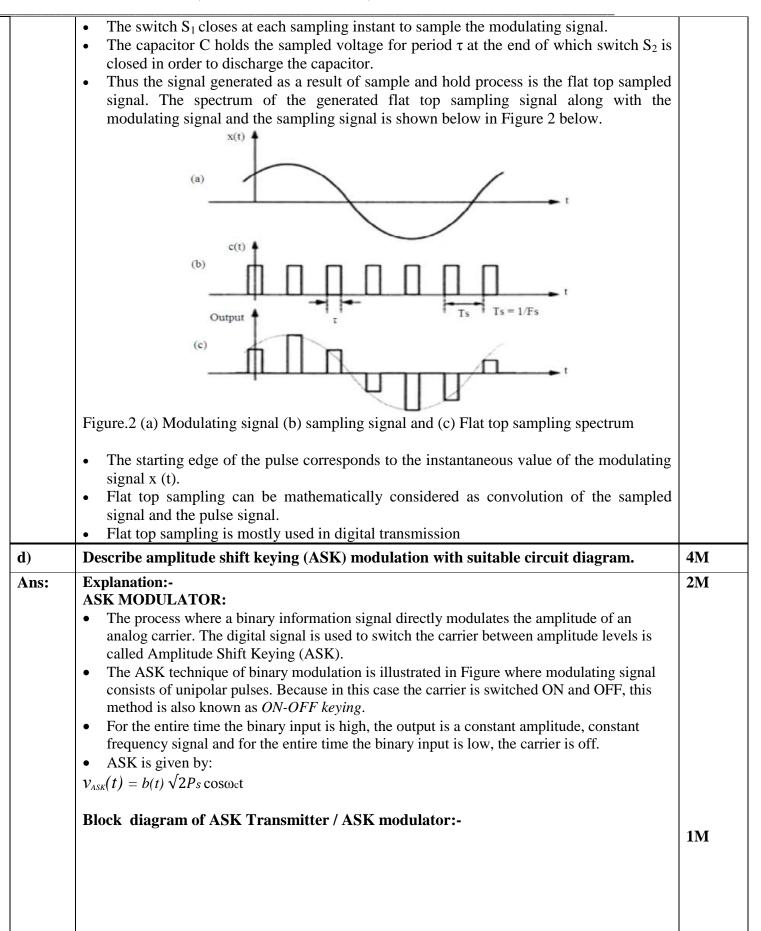


Ans:	Advantages of Digital Communication : (any 2)	1M
	1. High noise interference tolerance due to digital nature of the signal.	each
	2. With channel coding, error detection and correction at receiver is possible.	
	3. It provides us added security to our information signal i.e. Data encryption is possible for	
	greater security.	
	4. Cheaper due to advances in digital VLSI technology.	
	5. Digital information can be saved and retrieved when necessary.	
	6. Large data storage is possible.	
	Disadvantages of Digital Communication : (any 2)	
	1. Large System Bandwidth: - Digital transmission requires a large system bandwidth to	
	communicate the same information in a digital format as compared to analog format.	
	2. High power consumption (Due to various stages of conversion).	
	3. Needs synchronization	
	4.Sampling Error.	
b)	Draw the block diagram of DM transmitter.Explain each block in detail.	4 M
ns:	Block diagram of DM transmitter:-	2M
	Γ	
	Analog Semple HAAA + Delta PCM	
	Sampling pulse	
	Digital-to-analog converter	
	(DAC)	
	Updown (JD	
	Clock counter und 1=up 0=down	
		2M
	Explanation:-	Z 1 VI
	Sample and Hold:-	
	The input analog is sampled and converted to PAM signal, which is compared with the	
	output of the DAC. The output of the DAC is a voltage equal to the regenerated magnitude	
	of the previous sample, which was stored in the up-down counter as a binary number.	
	Up-down counter:-	
	The up-down counter is incremented or decremented depending on whether the previous	
	sample is larger or smaller than the current sample.	
	The up-down counter is clocked at a rate equal to the sample rate. Therefore the up-down	
	counter is updated after each comparison.	
	Initially the up-down counter is zeroed and DAC output is 0v.	
	Initially the up-down counter is zeroed and DAC output is 0v.	
	Initially the up-down counter is zeroed and DAC output is 0v. The first sample is taken and converted to a PAM signal, and compared with zero volts.	
	Initially the up-down counter is zeroed and DAC output is 0v. The first sample is taken and converted to a PAM signal, and compared with zero volts. The output of the comparator is a logic 1 condition (+v), indicating that the current sample	

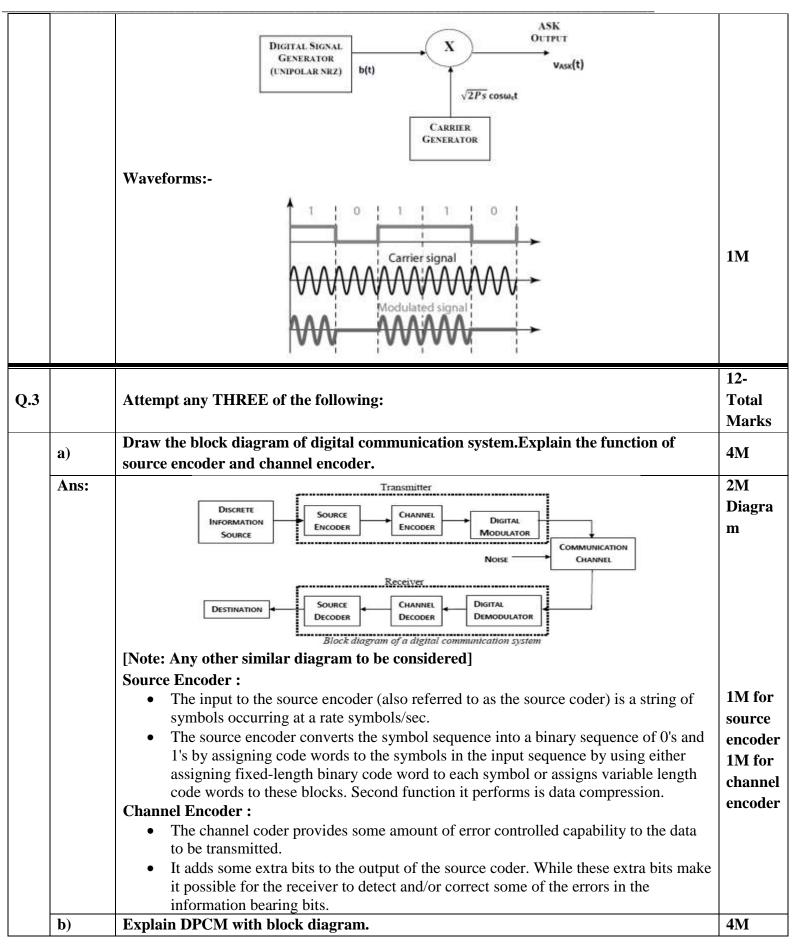














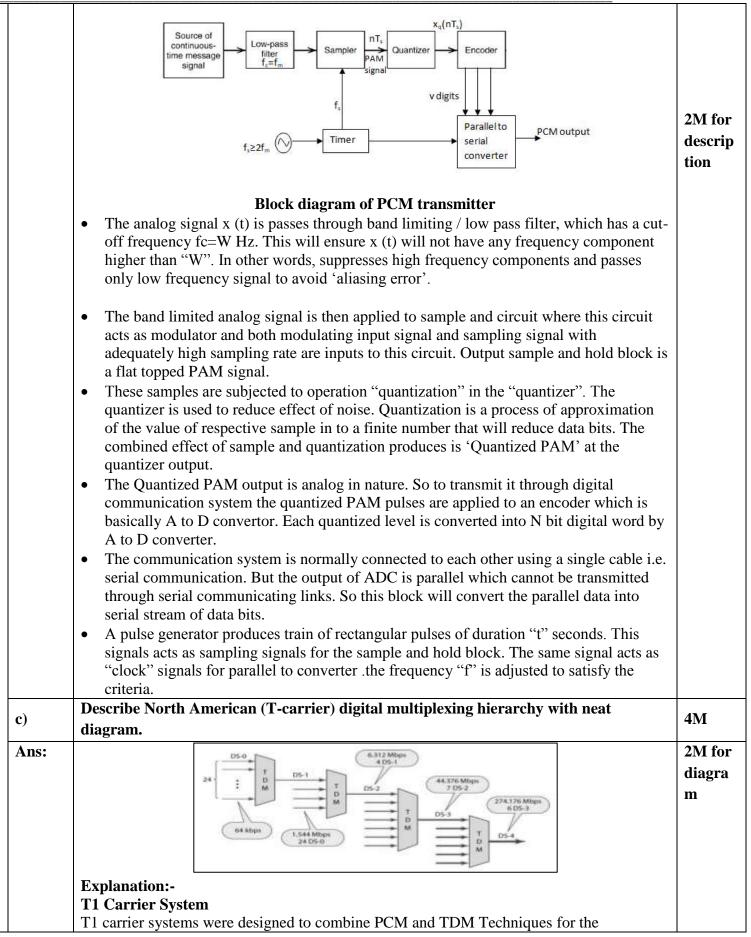
Ans:		Difference amplifier S	ampling signal	TRACT ENTRY	2M		
			Sample Cluantizer So(t) Encode	Parallel DPCM	diagra m		
	 Image: Accumulator + Predictor +						
				oximated signal x^(t)by passing			
			rough a predictor and accumulate				
c)			IA and CDMA (any four point		4M 1 mark		
Ans:	Sr. No.	Parameter	TDMA	CDMA	1 mark		
	1.	Definition	Entire bandwidth is shared among different subscribers at Fixed predetermined or dynamically assigned time intervals/slots.	Entire bandwidth is shared among different users by assigning unique codes.	Each point (Any 4		
	2.	Bandwidth Available	Time sharing of satellite transponder takes place	Sharing of bandwidth and time both takes place	point s)		
	3.	Synchronization	Synchronization is essential	Synchronization is not necessary			
	4.	Interference	Due to incorrect synchronization there can be interference between the adjacent time slots.	Both type of interference will be present			
	5.	Guard bands	Guard times between adjacent timeslots are necessary.	Guard bands and Guard times both are necessary			
	6.	Active terminals	Terminals are active in their specified slot on same frequency	All terminals active on same frequency			
	7.	Signal separation	Synchronization in time	Code separation			
	8.	Near Far Problem	No	Yes			
	9.	Handoff	Hard handoff	Soft handoff			
	10.	Application	Advanced mobile phone, system(AMPS), Cordless telephone	IS95 Wide band, CDMA 2000,2.5G and 3G			



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Sr.	T	Compare FDM & TDM systems (any four points).					
No.	FDM	TDM	1M for each				
1	Divides the channel into the two or more frequency ranges that do not overlap.	Divides and allocates certain Time periods to each channel.	point				
2	Code word is not required	No coding					
3	Needs guard bands	Needs guard time					
4	Problem of crosstalk	No problem of crosstalk					
Attem	pt any THREE of the following :		12- Total Marks				
State the Shannon Hartley's theorem for channel capacity. Explain the effect of S/N ratio and bandwidth on channel capacity.							
S/N ratio and bandwidth on channel capacity.In information theory, the Shannon–Hartley theorem tells the maximum rate at whichinformation can be transmitted over a communications channel of a specified bandwidth inthe presence of noise.According to Shannon, the bandwidth of the channel and signal energy and noise energy are related by the formula $C = W \log_2(1 + S/N)$ whereC is channel capacity in bits per second (bps)W is bandwidth of the channel in HzS/N is the signal-to-noise power ratio (SNR). SNR generally is measured in dB using the formula(S/N) dB = 10 log(Signal power / Noise power)Effect of S/N on Channel Capacity C:• If the communication channel is noiseless then N = 0. Therefore, S/N $\rightarrow \infty$ and so C also will tend to ∞ . Thus the noiseless channel will have an infinite capacity.Effect of Bandwidth B on Channel Capacity C:• Consider that some white Gaussian noise is present. Hence (S/N) is not infinite as N $\neq 0$. Now as the bandwidth approaches infinity, the channel capacity C does not become infinite because, N = η B will also increase with the bandwidth B. This will reduce the value of S/N							
Descri	Source of continuous- time message filter Sampler	Quantizer Encoder PCM signal applied to	4M 2M for block				
		R	diagra m				
	2 3 4 Attem State t S/N ra In info inform the pre Accord are rela C = W where C is ch W is b S/N is formul (S/N) 4 Effect • If the will ten Effect • Cons Now a becaus with in	or more frequency ranges that do not overlap.2Code word is not required3Needs guard bands4Problem of crosstalkAttempt any THREE of the following :State the Shannon Hartley's theorem for cl S/N ratio and bandwidth on channel capacidIn information theory, the Shannon-Hartley information can be transmitted over a commut the presence of noise.According to Shannon, the bandwidth of the c are related by the formula $C = W \log_2(1 + S/N)$ whereC is channel capacity in bits per second (bps) W is bandwidth of the channel in Hz S/N is the signal-to-noise power ratio (SNR). formula (S/N) dB = 10 log(Signal power / Noise pow Effect of S/N on Channel Capacity C: • If the communication channel is noiseless the will tend to ∞ . Thus the noiseless channel will Effect of Bandwidth B on Channel Capacit 	or more frequency ranges that do not overlap.Time periods to each channel.2Code word is not requiredNo coding3Needs guard bandsNeeds guard time4Problem of crosstalkNo problem of crosstalkAttempt any THREE of the following :State the Shannon Hartley's theorem for channel capacity. Explain the effect of S/N ratio and bandwidth on channel capacity.In information theory, the Shannon-Hartley theorem tells the maximum rate at which information can be transmitted over a communications channel of a specified bandwidth in the presence of noise.According to Shannon, the bandwidth of the channel and signal energy and noise energy are related by the formulaC = W log_(1 + S/N) whereC is channel capacity in bits per second (bps) W is bandwidth of the channel in Hz S/N is the signal-to-noise power ratio (SNR). SNR generally is measured in dB using the formula (S/N) dB = 10 log(Signal power / Noise power)Effect of S/N on Channel Capacity C: • If the communication channel is noiseless then N = 0. Therefore, S/N $\rightarrow \infty$ and so C also will tend to ∞ . Thus the noiseless channel will have an infinite capacity.Effect of Bandwidth B on Channel Capacity C: • Consider that some white Gaussian noise is present. Hence (S/N) is not infinite as N ≠ 0. Now as the bandwidth approaches infinity, the channel capacity C does not become infinite because, N =ηB will also increase with the bandwidth B. This will reduce the value of S/N with increase in B, assuming the signal power S to be constant.Describe PCM transmitter with block diagram.Impression of the signal power S to be constant.Describe PCM				







	Encoded voice band telepho a T1 carrier is 1.544 Mbps. All 24 DS-0 channels comb Called DS-1. Therefore T1	ined has a data	rate of 1.544M	bps, this digital s		2M for explana tion
	Service	Line	Rate (Mbps)	Voice Channels		
	DS-1	T-1	1.544	24		
	DS-2 DS-3	T-2 T-3	6.312	96		
	DS-4	T-4	274.176	4032		
<u>d)</u>	 DS and T Line rates T2 Carrier System T2 carriers time division metallic cable. T3 Carrier system T3 carriers Time division metallic cable. T4 Carrier System T4 carrier System T4 carriers time division metallic cable. T4 Carrier System T4 carriers time division metallic cable. T4 Carrier System T5 Carrier System T5 carriers time division metallic cable. T5 carriers time division metallic cable. T5 carrier System T	nultiplex 672 64 The transmission ultiplex 4032 64 ble upto 500 mil ultiplex 8064 64 single coaxial o	sted pair copper -kbps voice or on rate is 44.736 4-kbps voice or le. The transmis 4Kbps voice or cable.	wire up to 500 r data channels for 5 Mbps. data channels for sion rate is very data channels and	niles over a transmission transmitting high i.e. d transmit	4M
Ans:		nary modulation v band modulation PSK	Secondary mod Spread spectrum			2M for diagra
	In direct sequence, the serial binary data is mixed with a higher frequency pseudorandom binary code at a faster rate and the result is used to phase-modulate a carrier. The information signal undergoes primary modulation by PSK, FSK or other narrow band					
	modulation and secondary m obtained by multiplying the	odulation with s primary modulat vith commercial	pread spectrum ted signal and th radio, there are o	modulation. Sprea e square wave, ca cases where sprea	ad spectra are lled the PN d modulation is	tion
	sequence. Contrariwise, as w applied to the data first, and afterwards. The figure below parity. le of spread spectrun modulation.	is an example I	Hamming code	for the data 10	l0 with odd	
e)	applied to the data first, and a afterwards. The figure below parity. le of spread spectrum	is an example I n modulation and	Hamming code d demodulation	for the data 10 using PSK for pri	l0 with odd	4M



		T										
		Substitute, n=4	in the a	above ma								parity
					$\Rightarrow 2k \ge 4+$		_					bits.
						$5+k \Rightarrow 2k$						
		The minimum							,		1 I	
		bits p_1 , p_2 , and p_3 . Therefore, the number of bits in Hamming code will be 7, since there are 4 bits in binary code and 3 parity bits. We have to place the parity bits and bits of binary code in the Hamming code as shown below.										
	•										1.4	3
		 For a 4-bit code there are 3 parity bits p1, p2 and p3 at location 1, 2 and 4 resp. So, the code will be: "p1 p2 n1 p3 n2 n3 n4" where, n1, n2, n3, n4 are bits of the 										marks
		_							$1, n_2, n_3$	3, n	14 are bits of th	for
				-	are parity b							calculat
			re, the c	ode for e	ven parity	is calcu	lated	as below	:			ing
		statement		257	Bits	1		2	- 1			hammi
		Odd Parity fo	r dits 1,	3,3,7	P1	nl	ní		n4			
					0	1	0		0			ng code
			- h:+- 0 /		0	1	0		0			
		Odd parity for	r Dits 2,.	5,0,7	P2	nl	<u>n</u> .		n4			
					1	1	1		0			
			. 1. : 4	7	1	1	1		0			
		Odd parity for	r dits 4,:	5,6,7	P3	n2 0	n.		$\frac{n4}{n}$			
					0	0	1		0 0			
									0			
		Therefore ODI	O parity	hamming	g code wil	l be 011	0010					
Q.5		Attempt any T	WO of	the follo	wing:							12-
												Total
												Marks
	(a)	A diamete memory loss courses has an almhabet of some such the								e 6M		
	(a)	A discrete memory less source has an alphabet of seven symbols with probabilities f its output given in the following table:								1 UDADIIIIIes 10		
				-	5	C		G	C		C	
		$\begin{array}{ c c c c c c c c } Symbol & S_0 & S_1 & S_2 & S_3 & S_4 & S_5 & S_6 \\ \hline Symbol & S_0 & S_1 & S_2 & S_1 & S_2 & S_1 & S_2 & S_1 & S_2 &$										
		Probability 0.25 0.25 0.125 0.125 0.125 0.0625 0.0625										
	Compute:											
		(i) Huffman code for the above source.										
		(ii) The coding efficiency of the designed Huffman code.										
	Ans:	The Huffman	code fo	r the sou	rce is:							2M
		Symbol Stage I Stage II Stage III Stage IV Stage V Stage VI										
			5ª	*0.25 (25 - 02	5	5 F	0.5	+ 0.5	1		
			S.	0.26	25 402	5- 10:	15	*0.25	105	Ľ		
			5,		125 40.2	1	RT	025	G	1		
			S	1	125 -0.1	~	and .	1	1.294			
				V	0	¢-1	1		1			
			S4 0	N	1257 *0.1	1						
			2	4	The ancie	ried bile on	the date	d path corres	in the second			
			90	0.0825		for symbol S			Acres of			

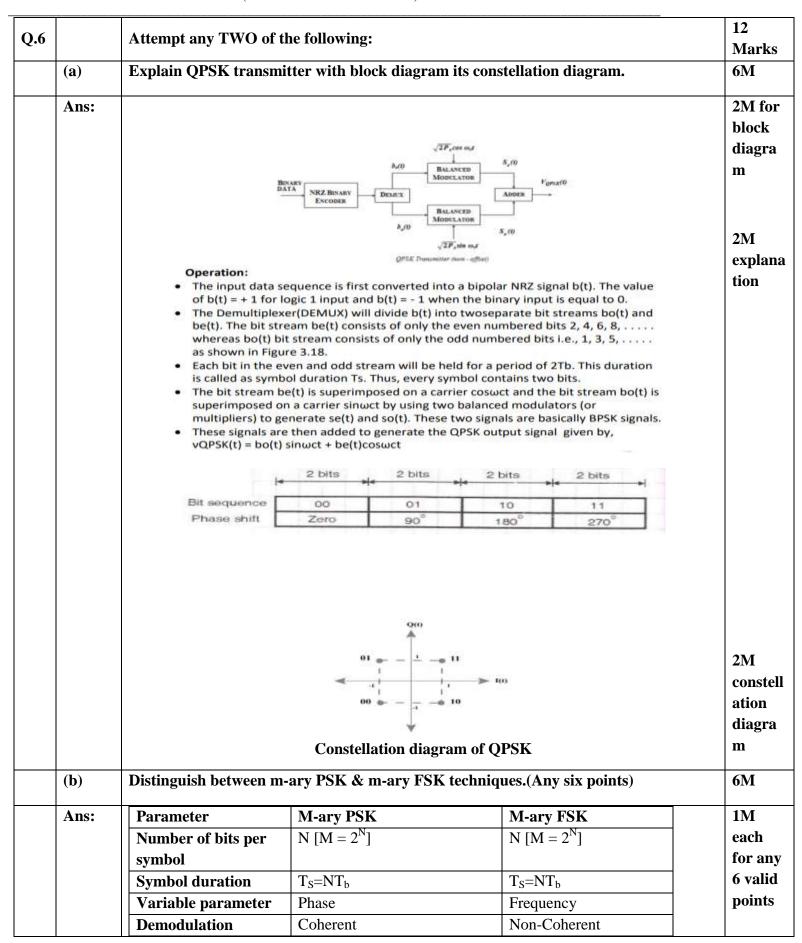


Symbol	Probability	Codeword	Codeword length	2M
Sa	0.25	10	2 bit	
Si	0.25	11	2 bit	
S2	0.125	001	3 bit	
Ss	0.125	010	3 bit	
54	0.125	011	3 bit	
S,	0.0625	0000	4 bit	
Se	0.0625	0001	4 bit	
k = 0				
	$P_s \times (\text{length of } s)$	ympon in chiev		
	able P. 2.7,3(b)		a second	
	$(0.25 \times 2) + (0.25 \times 2)$		(5×3)×3	
	+ (0.0625×4)			
	2.625 hits/symbo		and the second	
. The ave	rage information 6	per measage	and the second se	
= 1	$H = \sum_{i=0}^{n} p(x_i)$	log ₂ [1 / p (x _i	01	
л. н =	[0.25 log; (4)]	×2+[0.125]	og ₂ (8)]×3	
	+ [0.0625 log ₂ (16)]×2	and the second second	
= 1	[0.25×2×2]+	+ [0.125 × 3 × 3	3]+[0.0625×4×2]	
	2.625 bits/messa		- AND - REAL	
Code eff	ficiency $\eta = \frac{H}{L}$	$\times 100 = \frac{2.625}{2.625}$	× 100	
the second s	n = 100%		a strange of the set	
		114 4 1	niques (any six points)	. 6M



11	Sr. No	Parameter	Binary ASK	Binary FSK	Binary PSK	each for any
	1.	Variable Characteristic	Amplitude	Frequency	Phase	6 valid points
	2.	Maximum bandwidth(Hz)	2f _b	5 f _b /3	2f _b	
	3.	Noise immunity	low	high	high	
-	4.	Error probability	high	low	low	
	5.	Performance in presence of noise	poor	Better than ASK	Better than FSK	
	6.	Complexity	Simple	Moderately complex	Very complex	
	7.	Bit rate	Suitable upto 100 bits/sec	Suitable upto about 1200 bits/sec	Suitable for high bit rates	
	8.	Detection method	Envelope	Envelope	Coherent	
			tion reduces slope ov tion". Justify the abov			oise 6M
			tep size is constant so s . These drawbacks can	-	-	







	Method			
	Bandwidth	2fb/N	2^{N+1} fb/N	
	Probability of Error	More than that in M-ary FSK	Less than that in M-ary PSK	
	Transmitted signal	$\sqrt{2f_s} \cos (\omega_c t + \phi_m)$ $\therefore \phi_m = (2m + 1) \pi/4$	$\sqrt{2P_s}\cos\omega_c t$	
(c)	Explain fast frequency advantage and disadva	hopping techniques with suitab ntages.	ble waveforms. State its	6M
Ans:	 symbol. The hop rate is hit For each symbol symbol such that A jammer cannot than one carrier f 	MESK symbol 01 11 11 10 00 10 03 001100100100100100100100100100100100100	ate is equal to hop rate eral frequencies changes for one without is transmitted using more	2M explana tion 2M diagra m
	 3. Shorter time for a 4. Robust technolog Disadvantages : 1. Bandwidth require 2. Lower Coverage 	nsmission as only transmitter and acquisition gy rement is more [GHz] range due to high SNR requirem pensive digital frequency synthes		1M each for any one adavan tage and disadva ntage